



Intelligent Resource Allocation in Wireless Networks with Network Coding



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Selective Channelization

- Network Coding: Combine Packets Before Forwarding
 - Approaches multicast capacity of networks
 - Per-hop encoding and decoding employed
- Scenario: Congestion of multicast traffic
- Approach: Use TDMA MAC
 - Let network coding repair characteristics address low congestion
 - Channelize hyperarcs facing performance degradation
- Shortcomings:
 - Does not consider other characteristics of congested wireless networks - backoff, delays, queue occupancy, etc.
 - Does not take neighborhood into account before channelization

Cooperative Channelization

- Observation: Congestion arises out of competition for medium access in a neighborhood
- Define *health* metric using following components
 - Expected No. of Transmissions Due to Collisions
 - Expected No. of Media Access Attempts for Transmission
 - Network Coding Opportunity
 - Impact on Next-Hop Downstream Hyperarcs
- Channelization Scheme
 - A hyperarc is considered sick if its health < Threshold
 - Channelize sickest hyperarc in a 2-hop neighborhood

Expected No. of Transmissions Due to Collisions

Probability of node u receiving a packet from node v

$$p_u^v = (1 - h'_u \cdot q_u) \cdot \prod_w (1 - h'_w \cdot q_w), w \in \gamma_1(u) - \{v\}$$

$$h'_u = \begin{cases} h_u, & \text{if } h_u < 1 \\ 1, & \text{otherwise} \end{cases}$$

Expected No. of Transmissions Due To Collisions

$$X_v^i = T(|\eta_v^i|, |\eta_v^i|, \text{avg}_{u \in \eta_v^i} (p_u^i))_{nc} / |\eta_v^i|$$

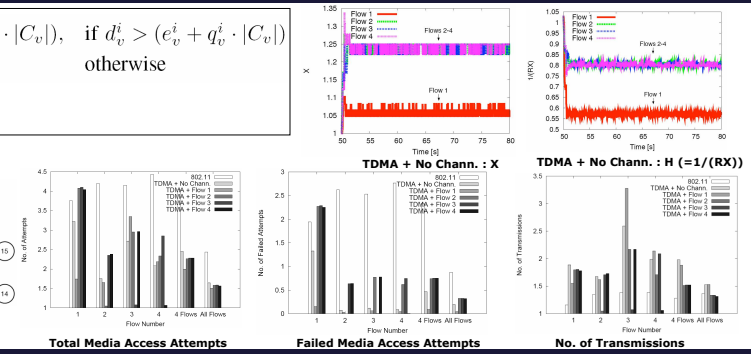
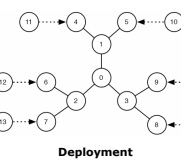
$$t_u = (1 - h'_u \cdot q_u) \cdot \prod_w (1 - h'_w \cdot q_w), w \in \gamma_1(u)$$

$$p_u^v = t_u / (1 - h'_v \cdot q_v)$$

Expected No. of Media Access Attempts For Transmission

$$R_v^i = \begin{cases} d_v^i / (e_v^i + q_v^i \cdot |C_v^i|), & \text{if } d_v^i > (e_v^i + q_v^i \cdot |C_v^i|) \\ 1, & \text{otherwise} \end{cases}$$

$$H_v^i = 1 / (R_v^i \cdot X_v^i)$$



Network Coding Opportunity

- Packets are never using XOR
- No more than 1 packet can be decoded out of a transmission
- No more than number of destinations packets can be coded
- X provides performance estimate of network coding under current conditions
- G provides performance estimate of network coding under varying conditions

$$G_v^i = 1 + |\eta_v^i| / |\gamma_1|_{max}$$

η_v^i - Set of destinations of the hyperarc
 $|\gamma_1|_{max}$ - Maximum number of 1-hop neighbors

$$H_v^i = G_v^i / (R_v^i \cdot X_v^i)$$

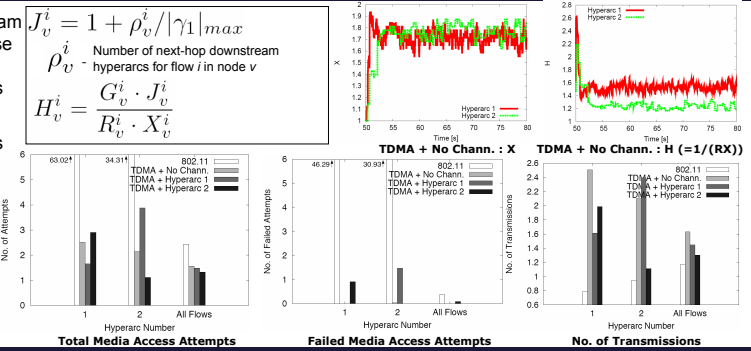
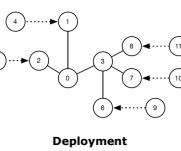
Impact on Next-Hop Downstream Hyperarcs

- Channelizing upstream hyperarcs may increase congestion on all downstream hyperarcs
- Channelize downstream hyperarcs first

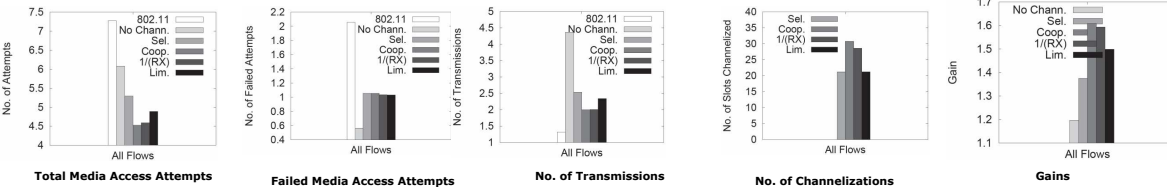
$$J_v^i = 1 + \rho_v^i / |\gamma_1|_{max}$$

ρ_v^i - Number of next-hop downstream hyperarcs for flow i in node v

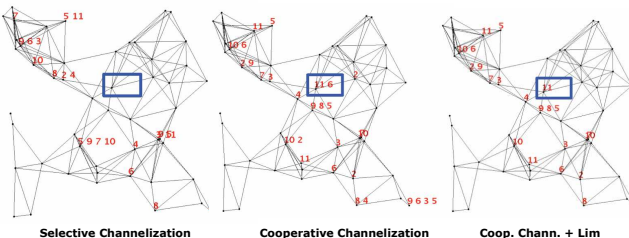
$$H_v^i = \frac{G_v^i \cdot J_v^i}{R_v^i \cdot X_v^i}$$



Multi-Hop Results



- Improvement over 802.11
 - TDMA + No Chann.: 20%
 - Selective Channelization: 38%
 - Cooperative Channelization: 61%
 - 1.5x more gain than Sel. Chann.



Varying Transmission Probability

